### **Parallel research of numerous eigenvalues in ABINIT**

# **François Alouges**<sup>(a)</sup>, **Christophe Audouze**<sup>(b)</sup> and , **Christophe Audouze**<sup>(b)</sup> and<br>**les Zérah**<sup>(b)</sup><br>- Laboratoire de Mathématiques<br>DAM, lle de France  $\boldsymbol{\mathsf{Gilles\text{ } Z}$ érah ${}^{(b)}$





### Framework of the problem

Calculation of the ground-state energy  $E^{KS}(\rho)$ , for a fixed geometry.

$$
\left\{\begin{array}{l} \bar{H}\varphi_i=\varepsilon_i\varphi_i,\ \forall i=1,\ldots,N,\\ \int_{\rm I\!R^3}\varphi_i\varphi_j^*=\delta_{ij}. \end{array}\right.
$$

### Framework of the problem

Calculation of the ground-state energy  $E^{KS}(\rho)$ , for a fixed geometry.

$$
\begin{cases} \bar{H}\varphi_i = \varepsilon_i\varphi_i, \ \forall i = 1, \dots, N, \\ \int_{\mathbb{R}^3} \varphi_i\varphi_j^* = \delta_{ij}. \end{cases}
$$

 Numerical resolution: discretization in a finite dimension basis  $\{ \chi_{j} \}_{j=1}$   $_M$  (Ga  $\boldsymbol{M}$  (Galerkin method)<br>eigenvalue problem generalized eigenvalue problem

$$
\varphi_i = \sum_{j=1}^{M} c_{ji} \chi_j \Rightarrow \tilde{H} \mathbf{c_i} = \varepsilon_i S \mathbf{c_i}, \forall i = 1, ..., N
$$

 $\sim$   $\sim$   $\sim$   $\sim$   $\sim$   $\sim$   $\sim$  $\mathcal{I}$ : discretized KS Hamiltonian;  $S$ : overlap matrix

### Numerical resolution

Direct minimization of  $E^{KS}(\rho),$  by nonlinear conjugate gradient (NLCG) algorithm

$$
\min \left\{ \sum_{i=1}^{N} (\tilde{H}\psi_i, \psi_i); (\psi_i, \psi_j) = \delta_{ij} \right\}
$$
  

$$
\psi_i
$$
: wave vectors; { $\varphi_i$ }: KS orbitals

At SCF convergence:  $pan(\psi_1, \ldots, \psi_N) = Span(\varphi_1, \ldots, \varphi_N)$ 

### Numerical resolution

Direct minimization of  $E^{KS}(\rho),$  by nonlinear conjugate gradient (NLCG) algorithm

$$
\min \left\{ \sum_{i=1}^{N} (\tilde{H}\psi_i, \psi_i); (\psi_i, \psi_j) = \delta_{ij} \right\}
$$
  

$$
\psi_i
$$
: wave vectors; { $\varphi_i$ }: KS orbitals

At SCF convergence:  $pan(\psi_1,\ldots,\psi_N)=Span(\varphi_1,\ldots,\varphi_N)$ :omnlexity (orthodonalization of the

Complexity (orthogonalization of the  $\psi_i$ ):  $\mathcal{O}(MN^2)$  or  $\sim$  0. . (matrix/vectors products in  $\mathcal{O}(N^2 \log N)$ )

Specificities of the ABINIT code and difficulties:

- for each SCF loop, calculation of pseudo-eigenvectors (i.e. not converged)
- : full symmetric matrix, of very large dimension and not explicitly known
- the eigenvalues calculation is intrinsically not parallel

Specificities of the ABINIT code and difficulties:

- for each SCF loop, calculation of pseudo-eigenvectors (i.e. not converged)
- : full symmetric matrix, of very large dimension and not explicitly known
- the eigenvalues calculation is intrinsically not parallel

Goal: propose parallel methods to realize calculations on more complex molecular systems.

# New algorithm (1)

calculation of exact eigenvectors instead of pseudo-eigenvectors  $\Rightarrow$  orthogonality of the  $\psi_i$  ( $\tilde{H}$ symmetric) such that

$$
\min \left\{ \sum_{i=1}^{N} ((\tilde{H} - \lambda_i I)^2 \psi_i, \psi_i) \right\}
$$

# New algorithm (1)

calculation of exact eigenvectors instead of pseudo-eigenvectors  $\Rightarrow$  orthogonality of the  $\psi_i$  ( $H$ symmetric) such that

$$
\min \left\{ \sum_{i=1}^{N} ((\tilde{H} - \lambda_i I)^2 \psi_i, \psi_i) \right\}
$$

spectral partitionning of an energy band  $\lceil a \rceil$  $\left[ a,b\right]$ iza bounds to be found (i.e. spectrum localization).

# New algorithm (1)

יונים וואו לא מודע הוא מודע ה<br>לא מודע הוא מודע הו

calculation of exact eigenvectors instead of pseudo-eigenvectors  $\Rightarrow$  orthogonality of the  $\psi_i$  ( $H$ symmetric) such that

$$
\min \left\{ \sum_{i=1}^{N} ((\tilde{H} - \lambda_i I)^2 \psi_i, \psi_i) \right\}
$$

spectral partitionning of an energy band  $\lceil a \rceil$  $\left[ a,b\right]$ iza bounds to be found (i.e. spectrum localization).

Proposals for the choice of  $a$  and  $b$ :

- given by the code user
- Lanczos algorithm
	- estimate of the  $N^{th}$  eigenvalue of  $\bar{H}$

#### Spectral partitionning: details

Iterative algorithm with dynamic allocations of shifts  $\lambda_i$ 



#### Spectral partitionning: details

Iterative algorithm with dynamic allocations of shifts  $\lambda_i$  $\bullet$ 



 $> N$  eigenvalues in  $[a, b] \Rightarrow \alpha k$  research of igenvalues (3  $< \alpha < 4$ ). eigenvalues ( $3 < \alpha < 4)$ .

To take into account of multiple eigenvalues:

$$
\min \left\{ \sum_{i=1}^{N_n} \sum_{j=1}^n \left( (\tilde{H} - \lambda_i I)^2 \psi_{ij}, \psi_{ij} \right), (\psi_{ij}, \psi_{ik}) = \delta_{jk} \right\},\,
$$

and the contract of the contra where  $N_{\bm n} =$  $\frac{N}{\bm{n}} \$  $\cdot$  is the number of blocks of  $\boldsymbol{n}$  wave vectors  $(n \in \{1, 2, 3, 4\})$ .

To take into account of multiple eigenvalues:

$$
\min \left\{ \sum_{i=1}^{N_n} \sum_{j=1}^n \left( (\tilde{H} - \lambda_i I)^2 \psi_{ij}, \psi_{ij} \right), (\psi_{ij}, \psi_{ik}) = \delta_{jk} \right\},\
$$

vectors  $(\bm{n} \in \{1,2,3,4\}).$ Theoretical complexity: where  $N_{\bm n} =$  $\frac{N}{\bm{n}} \$  $\cdot$  is the number of blocks of  $\boldsymbol{n}$  wave

 sequentially. Theoretical complexity:  $N_{\bm n}\times M{\bm n}^2={\bm n}MN\Rightarrow{\cal O}(N^2)$ 

#### Results and conclusions:

**•** Principle of the algorithm without reorthog. is validated  $(H_2, Si)$  but ... real complexity:  $\alpha N^3 \rightarrow \beta N^2$  with  $\beta \gg \alpha$ <br>because: because:

#### Results and conclusions:

 $\overline{a}$ 

- **•** Principle of the algorithm without reorthog. is validated  $(H_2,Si)$  but ... real complexity:  $\alpha N^3\rightarrow \beta N^2$ <br>because:<br>calculation of exact eigenvalues by NLCG because:
- with  $\beta \gg \alpha$ <br>can be very calculation of exact eigenvalues by NLCG can be very hard. The difficulties are:

numerical: cond(
$$
\psi_k
$$
) = 
$$
\frac{\lambda_{\text{max}}}{\min_{j \neq k} |\lambda_j - \lambda_k|}
$$

 $\frac{1}{\lambda_j - \lambda_k}$ <br>iltonians physical: spectrum of Hamiltonians has not <sup>a</sup> regular distribution

#### Results and conclusions:

 $\overline{a}$ 

- **•** Principle of the algorithm without reorthog. is validated  $(H_2,Si)$  but ... real complexity:  $\alpha N^3\rightarrow \beta N^2$ <br>because:<br>calculation of exact eigenvalues by NLCG because:
- with  $\beta \gg \alpha$ <br>can be very calculation of exact eigenvalues by NLCG can be very hard. The difficulties are:

numerical: cond(
$$
\psi_k
$$
) = 
$$
\frac{\lambda_{\text{max}}}{\min_{j \neq k} |\lambda_j - \lambda_k|}
$$

 $\frac{1}{\lambda_j - \lambda_k}$ iltonians physical: spectrum of Hamiltonians has not <sup>a</sup> regular distribution

 $\overline{L}$  and  $\overline{L}$  and

and  $\rho$  are invariant under any unitary tranformation<br>roposal of a new algorithm (2) with calculations of<br>ters of eigenvalues. proposal of <sup>a</sup> new algorithm (2) with calculations of clusters of eigenvalues.

### New algorithm (2)

Calculation of eigensubspaces with local reorthog. by

$$
\begin{cases} \text{ . } \text{NLCG on } (\tilde{H} - \lambda_i I)^2 \\ \text{ . } \text{Lanczos on } (\tilde{H} - \lambda_i I)^{-1} \end{cases}
$$

## New algorithm (2)

- Calculation of eigensubspaces with local reorthog. by NLCG on  $(\tilde{H})$  $(\lambda_i I)^2$ Lanczos on  $(\tilde{H})$  $(\lambda_i I)^{-1}$
- Hierarchy of algorithms:



#### • Numerical property of NLCG for eigensubspaces calculation:

**Proposition 1.** Let  $A$  be a matrix of order  $M$  and  $\{u_j, \lambda_j\}$  its eigenstates such that  $|\lambda_1 - \mu| < \ldots < |\lambda_M - \mu|.$  Then the<br>eenvergence fector of the enproximated eigenenese of dimen convergence factor of the approximated eigenspace of dimension  $m$  is  ${\cal O}$  $\frac{\lambda}{\lambda_n}$  $m - \mu$  $m+1 - \mu$  $\left(\begin{matrix} \end{matrix}\right)$  (subspace iteration algorithm).

#### Numerical property of NLCG for eigensubspaces calculation:

**Proposition 1.** Let  $A$  be a matrix of order  $M$  and  $\{u_j, \lambda_j\}$  its eigenstates such that  $|\lambda_1 - \mu| < \ldots < |\lambda_M - \mu|.$  Then the<br>eenvergence fector of the enproximated eigenenese of dimen convergence factor of the approximated eigenspace of dimension  $m$  is  ${\cal O}$  $\frac{\lambda}{\lambda_n}$ ner  $m - \mu$  $m+1 - \mu$ (subspace iteration algorithm).<br>
f the closest cluster near  $\mu$  tl

if  $m_{\mu}$  = dimension of the closest cluster near  $\mu$  then:

- $m < m_\mu$ : very slow convergence
- $m = m_{\mu}$ : optimal convergence

ؚ<br>ا

 $m > m_{\mu}$ : it depends of the  $\exists$  of other clusters

#### • Exemple: A diagonal,  $dim(A) = 20$ , concentrated spectrum:

 $\mathsf{sp}(A) = \{0.99$  $(0, 1.02, 1.03, 1.04, 1.05, \ldots, 3.99)$  $\{4.02, 4.03, 4.04, 4.05\}$ 

#### • Exemple: A diagonal,  $dim(A) = 20$ , concentrated spectrum:

 $\mathsf{sp}(A) = \{0.99$  $(0, 1.02, 1.03, 1.04, 1.05, \ldots, 3.99)$  $\{4.02, 4.03, 4.04, 4.05\}$ 



$$
\mu = -1.0, \, \log(|\lambda_i^{(k)} - \lambda_i|) = f(k), \, m = 5 \text{ et } 6.
$$

Paris, may  $2004 - p.11/13$ 





 $=-1.0$  $\log(|\lambda_i^{(k)}|)$  $\lambda^{(k)} - \lambda_i | ) = f(k), \ m = 10$  et 11





 $=-1.0$  $\log(|\lambda_i^{(k)}|)$ 

 $\begin{align} \mathcal{L}^{(k)} - \lambda_i | & = f(k), \ m = 10 \text{ et } 11 \end{align}$ <br>b-clusters of eigenvalues given<br>artitionning of  $[a,b],$  related to convergence of sub-clusters of eigenvalues gives another criterium of partitionning of  $\left[ a,b\right]$ , related to physical properties. physical properties. In a few words:

(1) - Spectral slice of an energy band with bounds to be chosen

(2) - Filling of energy regions via calculations of eigensubspaces, with local reorthog. Two ways:

> $\ddot{\phantom{a}}$ (2.1) - Lanczos on  $(\tilde H$  $(\lambda_i I)^{-1}$  "Classical" partitionning. (2.2) - NLCG on  $(\tilde H$  $(\lambda_i I)^2$  Physical partitionning related to the clusters of eigenvalues.

In a few words:

(1) - Spectral slice of an energy band with bounds to be chosen

(2) - Filling of energy regions via calculations of eigensubspaces, with local reorthog. Two ways:

> $\ddot{\phantom{a}}$ (2.1) - Lanczos on  $(\tilde H$  $(\lambda_i I)^{-1}$  "Classical" partitionning. (2.2) - NLCG on  $(\tilde H$  $(\lambda_i I)^2$  Physical partitionning related to the clusters of eigenvalues.

 algorithms (numerical and theoretical studies). "Long distance orthogonality": theorems quantifying the speed of orthog. between 2 vectors for different